

## ABSTRACT

A data encryption method performed with ring arithmetic operations using a residue number multiplication process wherein a first conversion to a first basis is done using a mixed radix system and a second conversion to a second basis is done using a mixed radix system. In some embodiments, a modulus  $C$  is chosen of the form  $2^w - L$ , wherein  $C$  is a  $w$ -bit number and  $L$  is a low Hamming weight odd integer less than  $2^{(w-1)/2}$ . And in some of those embodiments, the residue mod  $C$  is calculated via several steps.  $P$  is split into 2  $w$ -bit words  $H_1$  and  $L_1$ .  $S_1$  is calculated as equal to  $L_1 + (H_1 2^{x_1}) + (H_1 2^{x_2}) + \dots + (H_1 2^{x_k}) + H_1$ .  $S_1$  is split into two  $w$ -bit words  $H_2$  and  $L_2$ .  $S_2$  is computed as being equal to  $L_2 + (H_2 2^{x_1}) + (H_2 2^{x_2}) + \dots + (H_2 2^{x_k}) + H_2$ .  $S_3$  is computed as being equal to  $S_2 + (2^{x_1} + \dots + 2^{x_k} + 1)$ . And the residue is determined by comparing  $S_3$  to  $2^w$ . If  $S_3 < 2^w$ , then the residue equals  $S_2$ . If  $S_3 \geq 2^w$ , then the residue equals  $S_3 - 2^w$ .